# Successive minimum spanning trees Svante Janson <br> (joint work with Gregory Sorkin) 

Consider the complete graph $K_{n}$ with edge costs that are i.i.d. random variables, with a uniform distribution $U(0,1)$ (or, alternatively, an exponential distribution $\operatorname{Exp}(1))$. A well-known problem is to find the minimum (cost) spanning tree $T_{1}$, and its cost $c\left(T_{1}\right)$. A famous result by Frieze [2] shows that as $n \rightarrow \infty, c\left(T_{1}\right)$ converges in probability to $\zeta(3)$. (In both the uniform and exponential cases.)

Suppose now that we want a second spanning tree $T_{2}$, edge-disjoint from the first, and that we select it in a greedy fashion by first finding the minimum spanning tree $T_{1}$, and then the minimum spanning tree $T_{2}$ using only the remaining edges. (I.e., the minimum spanning tree in $K_{n} \backslash T_{1}$, meaning the graph with edge set $E\left(K_{n}\right) \backslash E\left(T_{1}\right)$.) We then continue and define $T_{3}$ as the minimum spanning tree in $K_{n} \backslash\left(T_{1} \cup T_{2}\right)$, and so on. We show that the costs $c\left(T_{2}\right), c\left(T_{3}\right), \ldots$ also converge in probability to some constants.

Theorem 1. For each $k \geq 1$, there exists a constant $\gamma_{k}$ such that, as $n \rightarrow \infty$, $c\left(T_{k}\right) \xrightarrow{\mathrm{p}} \gamma_{k}$ (for both uniform and exponential cost distributions).

The result extends easily to other distributions of the edge costs, by standard arguments, but we consider in here only the uniform and exponential cases.

By Frieze [2], $\gamma_{1}=\zeta(3)$. The constants $\gamma_{k}$ for larger $k$ are given by some expressions in the proof, but not in a form that is easily evaluated since they involve solutions of some non-linear functional equations (which furthermore involve a parameter). We can show the following bounds, which imply that $\gamma_{k}$ is roughly $2 k$ for large $k$ :

$$
\begin{equation*}
k^{2} \leq \sum_{i=1}^{k} \gamma_{i} \leq k^{2}+k, \quad k \geq 1 \tag{1}
\end{equation*}
$$

and

$$
\begin{equation*}
2 k-2 k^{1 / 2}<\gamma_{k}<2 k+2 k^{1 / 2}, \quad k \geq 1 \tag{2}
\end{equation*}
$$

A minor technical problem is that $T_{2}$ (and $T_{3}, \ldots$ ) does not always exist; it may happen that $T_{1}$ is a star and then $K_{n} \backslash T_{1}$ is disconnected. This happens only with a small probability, and w.h.p. (with high probability, i.e., with probability $1-o(1)$ as $n \rightarrow \infty), T_{k}$ is defined for every fixed $k$. However, we avoid this problem completely by modifying the model: we assume that we have a multigraph with an infinite number of copies of each edge in $K_{n}$, and that these have the costs given by the points in a Poisson process with intensity 1 on $[0, \infty)$. (The Poisson processes for different edges are, of course, independent.) Note that when finding $T_{1}$, we only care about the cheapest copy of each edge, and its cost has an $\operatorname{Exp}(1)$ distribution, so the problem for $T_{1}$ is the same as the original one. However, we now never run out of edges and we can define $T_{k}$ for all integers $k=1,2,3, \ldots$ Asymptotically, the three models are equivalent, and Theorem 1 holds for any of the models.

The multigraph model, moreover, is useful in our proofs because of the added independence.

Frieze [2] also proved that the expectation $\mathbb{E} c\left(T_{1}\right)$ converges to $\zeta(3)$. For the multigraph model just described, this too extends.

Theorem 2. For the multigraph model, $\mathbb{E} c\left(T_{k}\right) \rightarrow \gamma_{k}$ for each $k \geq 1$ as $n \rightarrow \infty$.
Remark 3. However, for the simple graph $K_{n}$ with, say, exponential costs, there is as said above a small but positive probability that $T_{k}$ does not exist for $k \geq 2$. Hence, either $\mathbb{E} c\left(T_{k}\right)$ is undefined for $k \geq 2$, or (better) we define $c\left(T_{k}\right)=\infty$ when $T_{k}$ does not exist, and then $\mathbb{E} c\left(T_{k}\right)=\infty$ for $k \geq 2$ and every $n$. Hence Theorem 2 does not hold for simple graphs, and the multigraph model is essential for studying the expectation.

Remark 4. Frieze and Johansson [3] recently considered a related problem, where instead of choosing spanning trees $T_{1}, T_{2}, \ldots$ greedily one by one, they choose $k$ edge-disjoint spanning trees with minimum total cost. It is easy to see, by small examples, that selecting $k$ spanning trees greedily one by one does not always give a set of $k$ edge-disjoint spanning trees with minimum cost, so the problems are different. We can also show that, at least for $k=2$, the two problems also asymptotically have different answers, in the sense that the limiting values of the minimum cost (which exist for both problems) are different.

The proofs are, as in many other previous papers on the random minimum spanning tree problem, based on Kruskal's algorithm which processes the edges in order of increasing cost and keeps the ones that join two different components in the forest obtained so far. (I.e., it keeps the edges that do not form a cycle together with some previously chosen edges.) The second minimum spanning tree can then be found by another application of the same algorithm to the remaining edges, and so on.

The results are proved by considering a random (multi)graph process, where copies of each edge $i j$ arrive as a Poisson process with intensity $1 / n$; an edge arriving at time $t$ has cost $t / n$. We let $G_{1}(t)$ be the multigraph formed by the edges that have arrived at time $t$. We run Kruskal's algorithm and let $F_{1}(t)$ be the forest formed by the edges selected up to time $t$ for the minimum spanning tree $T_{1}$. We let $G_{2}(t)$ be the multigraph consisting of the edges in $G_{1}(t) \backslash F_{1}(t)$, and let $F_{2}(t)$ be the forest formed by the edges selected up to time $t$ by Kruskal's algorithm applied to $G_{2}(t)$, and so on. We show, by induction in $k$, that each $G_{k}(t)$ is an example of an inhomogeneous random graph of the type studied in [1]; results from [1] thus yield results on the (asymptotic) structure of $G_{k}(t)$, in particular on the existence and size of a giant component, and these structural results are used to show the theorems above on the $\operatorname{cost} c\left(T_{k}\right)$.

## References

[1] Béla Bollobás, Svante Janson and Oliver Riordan, The phase transition in inhomogeneous random graphs. Random Struct. Alg. 31 (2007), 3-122.
[2] Alan M. Frieze, On the value of a random minimum spanning tree problem, Discrete Applied Mathematics 10 (1985), 47-56.
[3] Alan Frieze and Tony Johansson, On edge disjoint spanning trees in a randomly weighted complete graph. Preprint, 2015. arXiv:1505.03429

