

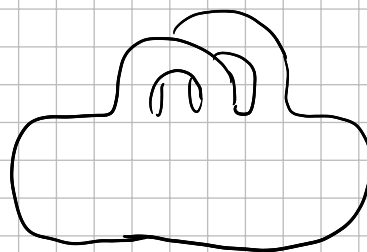
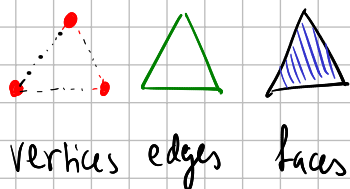
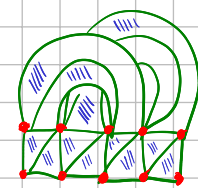
Knot polynomials

Before defining the Alexander polynomial.

How to find the genus of a surface with boundary

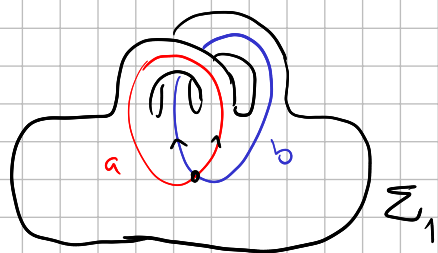
Method 1 Triangulate  $\Sigma$ , consider the Euler characteristic

$$\chi(\Sigma) = \#V - \#E + \#F = 2 - 2g - k$$


 $\cong$ 


$$10 - 23 + 12 = -1$$

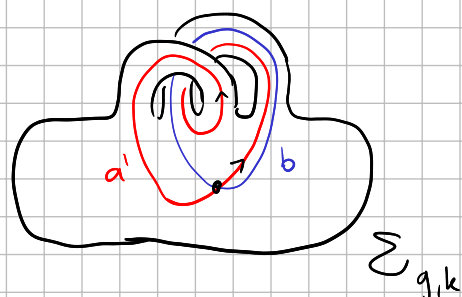
Method 2



$a, b$  form a basis of  $H_1(\Sigma_{g,k}) = \mathbb{Z}^{2g} \oplus \mathbb{Z}^{k-1}$

$g$ : genus

$k$ : nr of boundary components



$a' = a + b, b$  different basis

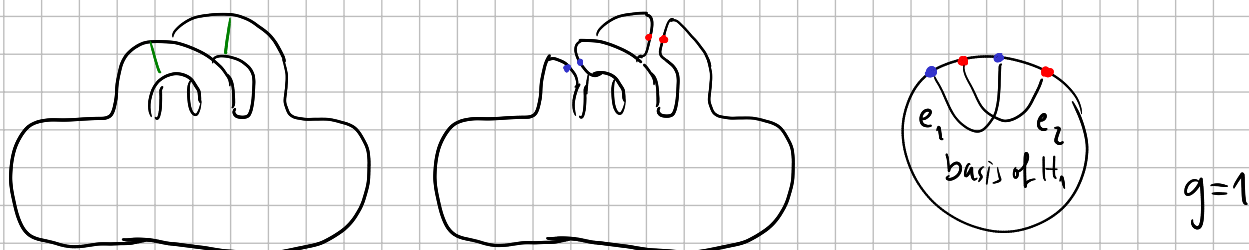
Our goal is to find the basis of simple closed curves as above for any  $\Sigma$ . But first, we need to find  $g$ .

(1) Cut  $\Sigma_{g,k}$  open along disjoint embedded arcs with boundary  $= \partial \Sigma$

Continue until we obtain a disc

(2) The number of arcs needed  $= 2g + (k-1)$

Ex

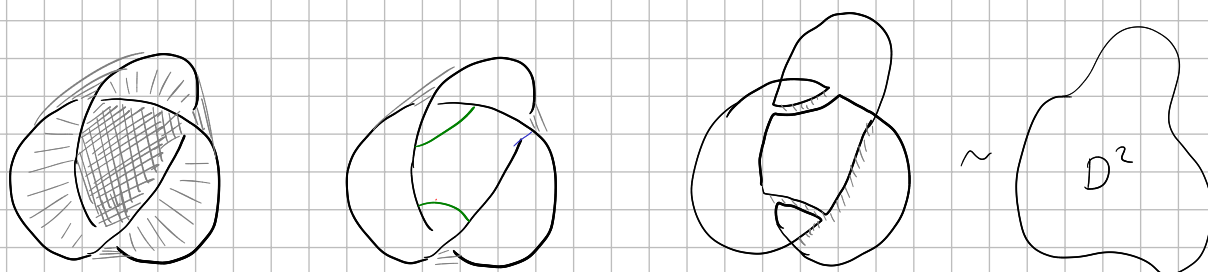


Rem • This reconstructs  $\Sigma_{g,k}$  as one-handle attachments on  $D^2$   
See Lecture 10.

- We obtain a basis of  $H_1$  by curves that intersect a unique arc transversely in a single point.

When  $g > 1$ , the handle decomposition is unfortunately not unique (there are handle-slides).

Ex



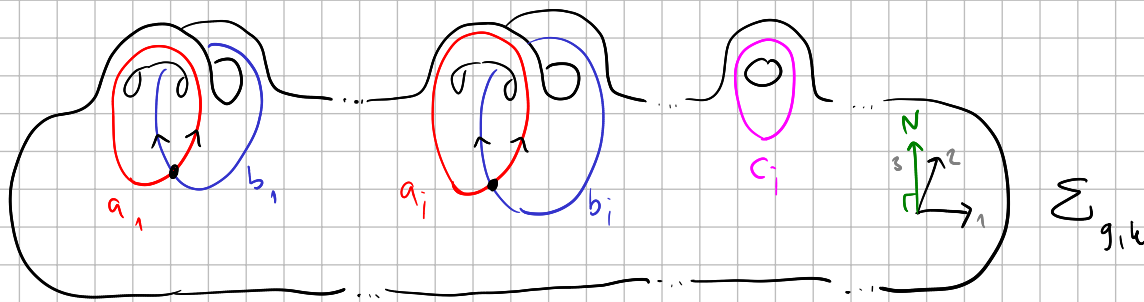
$$\Sigma_{g,k} \quad k=1: \quad 2g + (k-1) = 2 \Rightarrow g=1$$

$\Sigma_{g,k}$  has a basis of  $H_1(\Sigma_{g,k}) = \mathbb{Z}^{2g} \times \mathbb{Z}^{k-1}$  represented by simple closed curves

$$a_1, \dots, a_g, b_1, \dots, b_g, c_1, \dots, c_{k-1} \in H_1(\Sigma_{g,k})$$

where moreover  $x_i \cap x_j = \emptyset$  unless  $\bullet$   $x_i = x_j$

or  $\bullet$   $\{x_i, x_j\} = \{a_i, b_i\}$  single transverse intersection

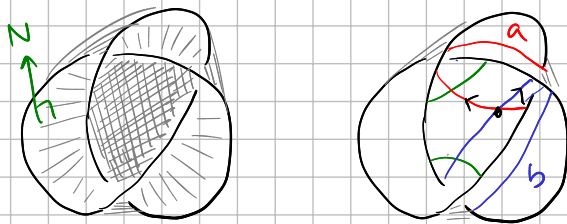


Given an orientation of  $\Sigma$ , we orient  $a_i$  &  $b_i$  so that the frame

$$\langle T_{a_i, n_{b_i}} a_i, T_{a_i, n_{b_i}} b_i \rangle = T_{a_i, n_{b_i}} \Sigma \text{ agrees with the orientation of } \Sigma.$$



Ex



For  $g > 1$ , the handle decomposition produced by the arcs need not give a basis of  $H_1$  with the correct intersection properties.

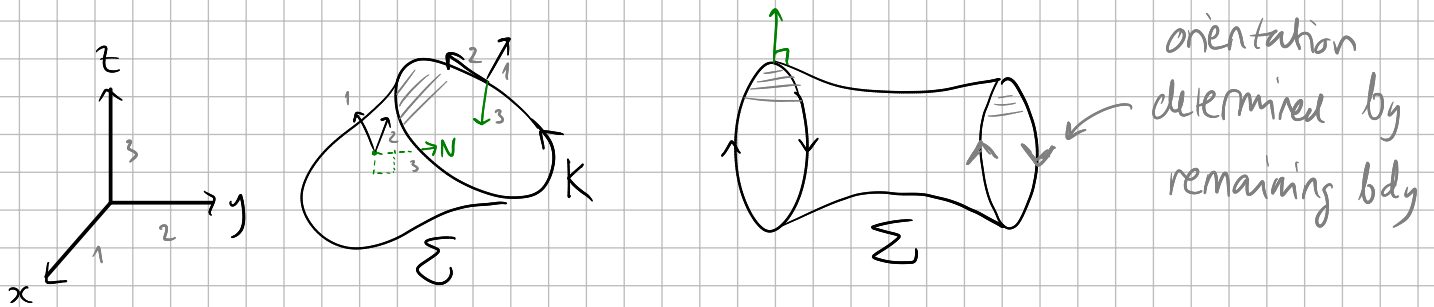
Fact There is a geometrically defined bilinear intersection form  
 $I: H_1(\Sigma) \otimes H_1(\Sigma) \rightarrow \mathbb{Z}$  ( $\Sigma \mapsto -\Sigma$  gives  $I \mapsto -I$ )  
 $I(a_i, b_i) = 1 = -I(b_i, a_i)$  opposite orient.

(orientability is crucial)

$\Rightarrow$  when  $g=0$ ,  $2g$  curves form a basis if they intersect like  $a$ 's &  $b$ 's

## The Alexander polynomial

Let  $K$  be an oriented knot/link, and  $\Sigma_{g, \pi_0(K)}$  a Seifert surface with orientation induced by  $K$ .



- If  $\Sigma$  is a Seifert surface ( $\Rightarrow$  conn. & orientable) then an orientation of a single component of  $K$  induces an orientation of the remaining components.
- The algorithm for constructing  $\Sigma$  from Lecture 16 depends on the orientation of  $K$ : it extends over the surface.

Choose a basis  $e_1 = a_1, e_2 = b_1, e_3 = a_2, e_4 = b_2, \dots, e_{2g} = b_g$

$e_{2g+1} = c_1, \dots, e_{2g+k-1} = c_{k-1}$

of simple closed curves with the above intersection properties (depends on the orientation of  $\Sigma$ , and hence on that of  $K$ ).

The Seifert matrix induced by  $\Sigma$  is

$V \in \text{Mat}_{2g+k-1, 2g+k-1}(\mathbb{Z})$  with entries

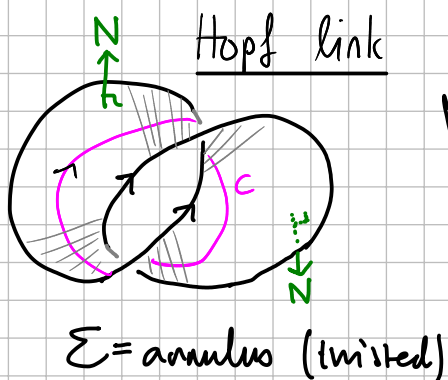
$$V_{j,i}^i = \text{lk}(e_i, e_j^\#)$$

row
pushoff of  $e_j$ 
along normal of  $\Sigma$ .

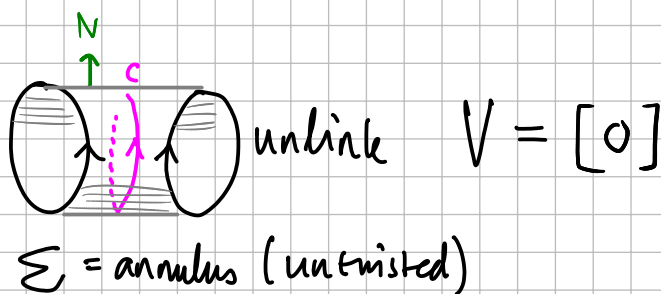
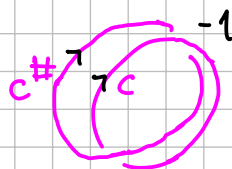
col.

$V$  is invariant / change of basis & certain stabilisations, but we will not focus  $V$  itself here.

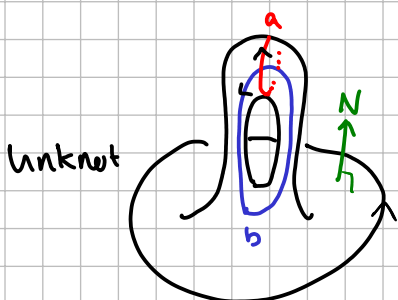
Ex



$$V = [\text{lk}(c, c^\#)] = [-1]$$



$$V = [0]$$



$$V = \begin{bmatrix} 0 & 0 \\ 1 & 0 \end{bmatrix}$$

$\text{lk}(a, b^\#)$

$\text{lk}(b, a^\#)$

The Alexander polynomial is the Laurent polynomial

$$\Delta_K(t) := t^N \cdot \det(tV - V^{\text{tr}}) \in \mathbb{Z}[t, t^{-1}]$$

$$N \in \mathbb{Z} \quad \text{s.t.} \quad \begin{cases} \bullet \Delta_K \text{ polynomial } & (\text{no negative } t \text{ powers}) \\ \bullet \Delta_K \text{ not divisible by } t & (\text{possibly zero}) \end{cases}$$

- ⚠
- there are different conventions for the normalisation
  - $\det(0 \times 0\text{-matrix}) = 1$

Ex

$$\Delta_{\emptyset\emptyset}(t) = 0, \quad \Delta_{\bigcirc}(t) = 1 - t,$$

$$\Delta_{\circlearrowleft}(t) = t^{-1} \cdot \det \begin{bmatrix} 0 & -1 \\ t & 0 \end{bmatrix} = 1.$$

Exercise 44

Show that (1)  $\pm t^{\deg \Delta} \Delta(t^{-1}) = \Delta(t)$

$$(2) \Delta_K(1) = \pm 1 \quad \text{if } |\pi_0(K)| = 1$$

Exercise 45

Calculate  $\Delta_{\bigcirc\bigcirc}(t) = 1 - t + t^2$

by using the above algorithm.

Thm

(Alexander, '28) The Alexander polynomial  $\Delta_K(t)$  is an invariant of oriented knots & links up to smooth isotopy.

# Skein relations

Some knot invariants can be determined via Skein-Relations.

Conway found the following version of the Alexander polynomial in the '60s.

$$\nabla_K(z) \in \mathbb{Z}[z^{\pm 1}], \quad \Delta_K(t) = t^{M/2} \cdot \nabla_K(t^{1/2} - t^{-1/2}) \quad M \text{ suitable}$$

- $\nabla_K$  is also an invariant of oriented knots

Conway showed

$$\nabla_{K_+}(z) - \nabla_{K_-}(z) = z \cdot \nabla_{K_0}(z)$$



Which together with  $\nabla_{\bigcirc}(z) = 1$  determines the invariant uniquely

Prop If a knot  $K$  has a diagram

then  $\nabla_K(t) = 0$ .

Proof

$K_+$	=		isotopic $\Rightarrow$	$\nabla_{K_+}(z) = \nabla_{K_-}(z)$
$K_-$	=			$\Rightarrow z \cdot \nabla_{K_0}(z) = 0$
$K_0$	=	= $K$		$\Rightarrow \nabla_{K_0}(z) = 0$ <span style="float: right;">□</span>

How to compute  $\Delta_K(t)$  using the Skein relation:

By induction, assume that  $\Delta_K(t)$  has been computed for all knots that admit a knot diagram with  $\leq n-1$  crossings.

Let  $K$  be a knot that admits a diagram with  $n$  crossings.

(1) For a suitable change  $X \rightsquigarrow Y$  at suitable subsets of the crossings, the new knot  $\tilde{K}$  is isotopic to  $O \cup O \cup \dots \cup O$  (unknot).



Proof Make  $z$ -coordinate increasing everywhere except along some small outer arc.  $\square$

(2) Change crossings of  $\tilde{K}$  one at a time to get back  $K$ .  
Solve the unknown term in the skein relation inductively.