

Problems with solutions

The numberings of results used refer to my notes, available on the course home page.

- (2.) Show that there exists a recursive function $h(x)$ such that for every $e \in \mathbb{N}$, if $\{e\}$ is total (i.e. defined for all $n \in \mathbb{N}$) and if $\text{ran}(\{e\})$ is infinite then $\{h(e)\}$ is an injective total function such that $\text{ran}(\{e\}) = \text{ran}(\{h(e)\})$.

Solution: Define $f(e, x) \simeq \{e\}(x)$.

Then by lemma 2.11, $f(e, x)$ is partial recursive

Now define $g(e, x) \approx$

$$\approx f\left(e, \left(\mu y \left[\forall u \leq x \forall v \leq x (u \neq v \rightarrow [(y)_u \neq (y)_v \wedge f(e, (y)_u) \neq f(e, (y)_v)]) \right] \right)\right)_x$$

Then g is partial recursive so by Corollary 2.12 there exists recursive function $h(x)$ such that $\{h(e)\}(x) \approx g(e, x)$ for all $e, x \in \mathbb{N}$.

It follows from the definition of g that if

$\{e\}$ is total and $\text{ran}(\{e\})$ is infinite

then $\{h(e)\}$ is total and $\text{ran}(\{h(e)\}) = \text{ran}(\{e\})$. $\square$

3. Show that the set

$A = \{e \in \mathbb{N} : W_e \text{ has at most 5 distinct elements}\}$
is not r.e.

Solution: First we show that $\bar{A} = \mathbb{N} - A$ is r.e. but not recursive.

We have $\bar{A} = \{e \in \mathbb{N} : W_e \text{ has at least 5 distinct elements}\}$ so

$$e \in \bar{A} \Leftrightarrow \exists y_1, \dots, y_5, z_1, \dots, z_5 \left[\underbrace{\bigwedge_{\substack{i, j \leq 5 \\ i \neq j}} y_i \neq y_j \wedge \bigwedge_{i=1}^5 T_i(e, y_i, z_i)}_{\text{recursive}} \right]$$

and therefore $\bar{A}$ is r.e.

By applying Rice's theorem one can show that $\bar{A}$ is not recursive (exercise!)

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By proposition 2.16 it follows that if $\bar{A}$ would also be r.e then $\bar{A}$ (and A) would be recursive, which contradicts what we have just proved. Hence $\bar{A}$ is not r.e. $\square$

4. Show that the relation R defined by $R(x, y) \Leftrightarrow \{x\}(y) \simeq 0$, is not recursive

Solution: Suppose that R is recursive.

Define $f(x, y) \simeq \begin{cases} 0 & \text{if } \{x\}(y) \downarrow \\ \uparrow & \text{if } \{x\}(y) \uparrow \end{cases}$

Then f is partial recursive (because

$f(x, y) \simeq Z(h(x, y))$ where Z is the zero function and $h(x, y) \simeq \{x\}(y)$, where h is partial recursive by lemma 2.11)

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so by Corollary 2.12 there is a recursive function $g(x)$ such that $\{g(x)\}(y) \simeq f(x, y)$ for all $x, y \in \mathbb{N}$. Now we get

$$K_0(x, y) \Leftrightarrow \{x\}(y) \downarrow \Leftrightarrow f(x, y) \simeq 0 \Leftrightarrow \{g(x)\}(y) \simeq 0$$

$\Leftrightarrow R(g(x), y)$ and since R and g are recursive it follows that K_0 is recursive, which contradicts proposition 2.20. Hence R is not recursive. $\square$

5. Let f be an injective recursive function and let $A = \text{ran}(f)$ (so it follows that A is infinite) and let $B = \{x \in \mathbb{N} : \text{there is } y > x \text{ such that } f(y) < f(x)\}$

(i) Show that B is r.e.

(ii) Show that if there is an infinite r.e. set $C \subseteq \mathbb{N}$ such that $B \cap C = \emptyset$ then A is recursive.

Solution: (i) Define $g(x) \simeq \mu y (y > x \wedge f(y) < f(x))$

Then g is partial recursive and $g(x) \downarrow$ if and only if $x \in B$, so $\text{dom}(g) = B$. Hence (by prop 2.15) B is r.e.

(ii) Suppose that $C \subseteq \mathbb{N}$ is infinite and r.e. and that $B \cap C = \emptyset$. Then there is (by prop 2.15) a recursive function h such that $\text{ran}(h) = C$.

Define $H(x) \simeq (\mu y (h(y) = (y)_0 \wedge (y)_0 > x))_0$

Then $H(x)$ is a total recursive function (because C is infinite), and $H(x) =$ the least n in C such that $n > x$.

It follows that there exists z such that $f(z) = x$ if and only if there exists $z < H(x)$ such that $f(z) = x$.

Then $x \in A \iff \exists z < H(x) (f(z) = x)$

and the right hand side is a recursive relation

Hence A is recursive.

